

Cyclability, Connectivity and Circumference

Anish Hebbar¹ and Niranjana Balachandran²

¹ Indian Institute of Science, Bangalore, India

² Indian Institute of Technology Bombay, India
anishhebbbar@iisc.ac.in niranj@iitb.ac.in

Abstract. In a graph G , a subset of vertices $S \subseteq V(G)$ is said to be cyclable if there is a cycle containing the vertices in some order. G is said to be k -cyclable if any subset of $k \geq 2$ vertices is cyclable. If any k ordered vertices are present in a common cycle in that order, then the graph is said to be k -ordered. We show that when $k \leq \sqrt{n+3}$, k -cyclable graphs also have circumference $c(G) \geq 2k$, and that this is best possible. Furthermore when $k \leq \frac{3n}{4} - 1$, $c(G) \geq k+2$, and for k -ordered graphs we show $c(G) \geq \min\{n, 2k\}$. We also generalize a result by Byer et al. [1] on the maximum number of edges in nonhamiltonian k -connected graphs and show that if G is a k -connected graph of order $n \geq 2(k^2 + k)$ with $|E(G)| > \binom{n-k}{2} + k^2$ then the graph is hamiltonian, and moreover the extremal graphs are unique.

Keywords: Cyclability · Connectivity · Circumference · Hamiltonicity

1 Introduction

We consider only finite, undirected, simple graphs throughout this paper. The vertex and edge sets of G will be denoted by $V(G)$ and $E(G)$ respectively, the graph complement by \overline{G} . The length of the longest cycle in the graph G , also known as the circumference, will be denoted by $c(G)$. The minimum degree, independence number and connectivity of a graph will denoted by $\delta(G)$, $\alpha(G)$ and $\kappa(G)$ respectively. We will also use $d_H(v)$ for the degree of v in H . The set of neighbours of a vertex $v \in V(G)$ will denoted by $N(v)$, and the closed neighbourhood of v , viz. $N(v) \cup \{v\}$ will be denoted by $N[v]$.

A subset $S \subseteq V(G)$ of vertices in a graph G is said to be cyclable if G has a cycle containing the vertices of S in some order. A graph G is said to be k -cyclable if any $k \geq 2$ vertices of G lie on a common cycle. Note that the problem of determining the hamiltonicity of a graph is a special case of cyclability, namely when $k = n$. Cyclability and connectivity are interlinked, as was shown by Dirac [2] who proved for every $k \geq 2$, k -connected graphs are also k -cyclable. In fact, for $k = 2$ connectivity and cyclability are equivalent, but in general for $k \geq 3$ it is not necessarily true that every k -cyclable graph is also k -connected, as can be seen by considering the graph $K_2 \vee 2K_k$ which has connectivity exactly 2 and is also k -cyclable. For a brief survey of results involving conditions for cycles to

contain a particular set, refer to [3].

There is a rich literature on conditions guaranteeing the presence of long cycles in graphs, the most classical one being that of Dirac [4] who showed that in 2-connected graphs, $c(G) \geq \min\{n, 2\delta(G)\}$. Moreover, k -connected graphs have circumference at least $\min\{n, 2k\}$ from an easy consequence of Menger's theorem, and this is tight. A famous result by Chvátal and Erdős [5] relates the connectivity and independence number of a graph to hamiltonicity, and says that if the connectivity of a graph G is at least its independence number, then the graph is hamiltonian. However, not much is known when the requirement of connectivity is weakened to cyclability. Bauer et al. [6] obtained lower bounds for the length of the longest cycle in 3-cyclable graphs in terms of the minimum degree and independence number, but not much else is known for k -cyclable graphs for arbitrary k .

Cyclability has also received interest from an algorithmic and complexity theoretic point of view as it is a 'hard' parameter that can be thought of as a more quantitative measure of hamiltonicity. Since the classical HAMILTONIAN CYCLE problem is NP-complete, the problem of determining whether a graph is k -cyclable (CYCLABILITY) is NP-complete as well. The problem of determining whether a given subset S of vertices is cyclable (TERMINAL CYCLABILITY) has been studied in the Parameterized Complexity framework (FPT) (parameterized by $|S|$) and the best known algorithm has running time $O(2^{|S|}n^{O(1)})$ [7]. For some special classes of graphs such as interval graphs and bipartite permutation graphs, Crespelle and Golovach [8] showed that both these problems can be solved in polynomial time. For $|S| = O((\log \log n)^{1/10})$, Kawarabayashi [9] obtained a polynomial time algorithm for TERMINAL CYCLABILITY.

Note that k -connectivity guarantees $c(G) \geq \min\{n, 2k\}$ and also ensures k -cyclability. Thus, a natural question to ask is whether the same bound on the circumference can be obtained when the connectivity criteria is weakened to cyclability. When $k = n - 1$, we would require any set of $n - 1$ vertices of G to lie on a common cycle. It turns out that in this case, it is not necessary that the graph is hamiltonian. Indeed, the existence of hypohamiltonian graphs [10] of order n is known for all $n \geq 18$. Our first result in this paper gives a similar circumference bound for a wide range of k :

Theorem 1. *Let G be a k -cyclable graph, where $2 \leq k \leq n$. Then,*

$$c(G) \geq \begin{cases} 2k & \text{if } k \leq \sqrt{n+3} \\ k+2 & \text{if } k \leq \frac{3n}{4} - 1 \end{cases}$$

Moreover, for $2 \leq k \leq \sqrt{n+3}$, this bound on the circumference is best possible.

Note that for $k \geq \frac{n}{2}$ it is still possible that one can have a bound of the form $c(G) \geq (1 + \gamma)k$ for some fixed positive constant $\gamma < 1$ as long as $k \neq n - O(1)$.

A related notion is the orderedness of a graph, a strong hamiltonian property that was first introduced by Ng and Schultz [11]. A graph G is said to be k -ordered if any sequence of distinct vertices $T = \{v_1, \dots, v_k\}$ are present in some common cycle, in that order. Note that k -ordered graphs are naturally also k -cyclable, and it is also easy to see that they are $(k - 1)$ -connected. For a comprehensive survey of results on k -ordered graphs, see [12]. We show that for k -orderedness, the same circumference bound as k -connectivity holds for all $2 \leq k \leq n$.

Theorem 2. *Let G be a k -ordered graph, $2 \leq k \leq n$. Then, $c(G) \geq \min\{n, 2k\}$.*

Our second pursuit in this paper is to obtain Turán-type results for the circumference of k -connected graphs, specifically the maximum number of edges in nonhamiltonian k -connected graphs. A classical result states that if G is a graph of order n with $|E(G)| > \binom{n-1}{2} + 1$, then G is hamiltonian. This was generalized by [1] for $k \leq 3$, where they showed that if G is k -connected and satisfies $|E(G)| > \binom{n-k}{2} + k^2$ with n sufficiently large, then the graph is hamiltonian and the extremal graphs are unique. We further generalize their result and extend it to any k satisfying $n \geq 2(k^2 + k)$.

Theorem 3. *Let G be a k connected graph of order $n \geq 2(k^2 + k)$. If $|E(G)| > \binom{n-k}{2} + k^2$, then G is hamiltonian. Moreover, the extremal graphs are unique.*

The rest of the paper is organized as follows. We lay out some preliminaries in the next section, and give the proofs of Theorems 1, 2, and 3 in the following section. We conclude with some remarks and open questions.

2 Preliminaries

When the underlying graph is clear, we will use δ, κ, α instead of $\delta(G), \kappa(G), \alpha(G)$ for brevity, and also omit the subscript in $d_H(v)$. We also use the following well-known lemma attributed to Dirac repeatedly throughout the paper, and provide an outline of the proof for completeness.

Lemma 4 ([2]). *Any k -connected graph G is k -cyclable. Moreover, it satisfies $c(G) \geq \min\{n, 2k\}$*

Proof Sketch. Suppose some subset S of vertices with $|S| = k$ was not fully contained in any cycle. Then, take a cycle C containing as many of the vertices of S as possible, and pick some $v \in S$ that is not in C . By Menger's theorem, we can choose k vertex-disjoint paths from v to the cycle C , and these endpoints divide C into k segments. Since there are strictly less than k vertices of S in C , one of the segments does not contain any vertex from S , and thus we can extend this segment with the 2 disjoint paths from v at the ends of the segment to obtain a cycle containing more vertices of S , contradiction.

Now consider the longest cycle C in G and suppose its length is strictly less than $\min\{n, 2k\}$. Pick some $v \in V(G)$ not in C , and by Menger's theorem there

are k vertex disjoint paths from v to C . By the pigeonhole principle, some two endpoints of these k paths must be adjacent on the cycle C , giving a contradiction as we can replace the edge between these endpoints with the 2 paths to obtain a longer cycle.

A famous result by Chvátal and Erdős states the following

Theorem 5 ([5]). *If in a graph G , $\alpha(G) \leq \kappa(G)$, then G is hamiltonian.*

A natural generalization of the above is to flip the condition $\alpha(G) \leq \kappa(G)$, and instead ask for lower bounds on the circumference of a graph G where $\alpha(G) \geq \kappa(G)$. Foquet and Jolivet [13] conjectured the following, which was later proven by Suil O, Douglas B. West and Hehui Wu.

Theorem 6 ([14]). *If G is a k -connected n -vertex graph with independence number α and $\alpha \geq k$, then G has a cycle of length at least $\frac{k(n+k-\alpha)}{\alpha}$.*

The following result by Dirac is well-known and was a precursor to a number of results involving the length of the longest cycle in a graph.

Theorem 7 ([4]). *If G is 2-connected and has minimum degree δ , $c(G) \geq \min\{2\delta, n\}$.*

Note that 2-connectivity is equivalent to 2-cyclability. Bauer et al. obtained a bound on the circumference of 3-cyclable graphs in terms on the minimum degree and independence number.

Theorem 8 ([6]). *If G is 3 cyclable, then*

$$c(G) \geq \min\{n, 3\delta - 3, n + \delta - \alpha\}.$$

Ng and Schultz studied a related hamiltonian property termed k -orderedness, and showed the following connectivity result. Once again, we include the proof for completeness.

Lemma 9 ([11]). *Let G be a k -ordered graph. Then, G is $(k - 1)$ -connected.*

Proof. If not, there exists a set S of $k - 2$ vertices whose removal disconnects G , breaking it into at least 2 components. Take 2 vertices u, v in different components, then any path from u to v must go through some vertex of S . Thus, let T consist of u, v and then the vertices of S , in that order. These vertices must appear in some cycle in that order, giving a contradiction.

We will also need the concept of graph closure introduced by Bondy and Chvátal. Define the closure of G , denoted $cl(G)$, to be the graph obtained by repeatedly joining any two nonadjacent vertices x, y that satisfy $d(x) + d(y) \geq n$ in G . They showed that $cl(G)$ is well-defined (independent of the order in which nonadjacent vertex pairs are considered), and that G is hamiltonian if and only if $cl(G)$ is also hamiltonian.

Lemma 10 ([15]). *Suppose $cl(G) = G$ for a nonhamiltonian graph G of order n . Then $d(x) + d(y) \leq n - 1$ for any pair $\{x, y\}$ of nonadjacent vertices.*

This was later generalized to obtain results for higher order connectivity, the bounds now also involving the independence number. We define

$$\sigma_k(G) = \min\left\{\sum_{i=1}^k d(x_i), \{x_1, \dots, x_k\} \text{ an independent set of size } k \text{ in } G\right\}$$

Note that $\sigma_1(G)$ simply corresponds to the minimum degree δ , and Ore's theorem [16] states that if $\sigma_2(G) \geq n$, then the graph is hamiltonian.

Theorem 11 ([17]). *Let G be a k connected graph of order n and independence number α . If $\sigma_{k+1}(G) \geq n + (k - 1)\alpha - (k - 1)$, then G is hamiltonian.*

3 Proofs of the Results

Proof of Theorem 1.

We will first prove the bound for the regime $2 \leq k \leq \sqrt{n+3}$. Consider any k -cyclable graph with $\alpha(G) \geq k$. Then, let S be a set of k independent vertices, and consider the cycle containing it. This gives us a cycle of length at least $2k$, as any 2 independent vertices are not adjacent to each other. Thus, we can assume $\alpha(G) \leq k - 1$. Let the connectivity of the graph be κ . Using Theorem 6, it suffices to show

$$\frac{\kappa(n + \kappa - \alpha)}{\alpha} \geq 2k \iff n \geq 2k\left(\frac{\alpha}{\kappa}\right) + (\alpha - \kappa)$$

As k -cyclable graphs are also 2-cyclable, and thus 2-connected, we must have $\kappa \geq 2$. Hence, it is sufficient if

$$n \geq 2k\left(\frac{k-1}{\kappa}\right) + k - 3$$

which is always true when

$$n \geq k^2 - 3 \iff k \leq \sqrt{n+3}$$

Note that if we only ask for an improvement of the form $c(G) \geq (1 + \gamma)k$ for some positive constant $\gamma < 1$, by the same argument as above we get $\alpha < \frac{(1+\gamma)k}{2}$, and it would suffice to have

$$n \geq \frac{(1 + \gamma)^2 k^2}{4} + \frac{(1 + \gamma)k}{2} - 2 \iff \frac{\sqrt{4n+9}}{1 + \gamma} \geq k$$

So the above argument only yields a linear improvement for k up to around $2\sqrt{n}$.

Now, suppose $2 \leq k \leq \frac{3n}{4} - 1$, and assume to the contrary that $c(G) < k + 2$. We must have $k \geq 3$ as 2-cyclable graphs are 2-connected and hence

have circumference at least 4 for $n \geq 4$. By Theorem 7, we must have $\delta \leq \frac{k+1}{2}$. Moreover, $\alpha \leq \frac{k+1}{2}$ as otherwise we could simply take a cycle containing α many independent vertices. Consider a vertex v with minimum degree δ , with neighbourhood $N(v)$ satisfying $|N(v)| = \delta$. Now, choose v and any $k-1$ vertices from $V \setminus N[v]$, which is possible as long as $k-1 \leq n-1-\delta$. Then, any cycle containing these vertices must also contain some 2 neighbours of v , giving $c(G) \geq k+2$, and we are done.

Thus, we must have $k + \delta > n$. Note that when $2 \leq k \leq \frac{3n}{4} - 1$, $n \geq k+2$ if $n \geq 4$. So, we must either have $3\delta - 3 \leq k+1$ or $n + \delta - \alpha \leq k+1$, otherwise we are done by Theorem 8.

The former inequality gives $\delta \leq \frac{k+4}{3}$, which gives

$$n < k + \delta \leq \frac{4k+4}{3} \implies \frac{3n-4}{4} < k$$

a contradiction. Hence, we must have $\delta \geq \frac{k+5}{3}$, $\alpha \leq \frac{k+1}{2}$ giving

$$k+1 \geq n + \delta - \alpha \geq n + \frac{k+5}{3} - \frac{k+1}{2} = n + \frac{7-k}{6}$$

or equivalently, $\frac{3n}{4} - 1 \geq k \geq \frac{6n+1}{7}$, which is again a contradiction.

We now prove an analogous bound for the circumference of k -ordered graphs.

Proof of Theorem 2.

We know that k -ordered graphs are also $k-1$ connected from Theorem 9, thus $\kappa \geq k-1$. We also must have $\alpha \leq k-1$, as otherwise we can simply take k independent vertices in any order to obtain a cycle of size at least $2k$, in which case we are done. Hence,

$$\kappa \geq k-1 \geq \alpha$$

so by Theorem 5, we have that G is hamiltonian, and thus we are done in this case as well.

In fact, it is not hard to see that the $\min\{n, 2k\}$ bound on the circumference is achieved for all $2 \leq k \leq n$. If $k > n/2$, simply consider the complete graph K_n which is clearly k -connected, k -ordered, k -cyclable and has circumference n . If $k \leq n/2$, consider the complete bipartite graph $G = K_{k, n-k} = (A, B, E)$, which is k -ordered, and hence k -cyclable. Indeed, take any sequence of k distinct vertices $T = (v_1, v_2, \dots, v_k)$. We construct a cycle containing T in that order as follows.

Let T_A be the set of vertices in T and A , with T_B being defined similarly. Then, for any $v \in T_A$, if the next vertex in the sequence T is in T_B , then simply follow the edge joining them. Otherwise, first follow an edge to a vertex in $B \setminus T_B$, and then back to the next vertex which must have been in T_A . Follow the same procedure for vertices in T_B . At the end, follow the edge joining the first and last vertex. We cannot run out of vertices as the number of extra vertices outside

T_A in A that are needed is at most $|T_B|$, and $|A| = k = |T_A| + |T_B|$. Similarly, $|B| = n - k \geq k = |T_A| + |T_B|$.

We now generalize a result by [1] on the maximal number of edges in a k -connected nonhamiltonian graph, for $k = 2, 3$. We will need the following short lemma which appears in [1].

Lemma 12 ([1]). *Let G be a nonhamiltonian, k -connected graph of order n . Then $k \leq \frac{n-1}{2}$ and $|E(\overline{G})| \geq \binom{k+1}{2} + (k-1)(n-k-1) - \sigma_{k+1}(G)$*

Proof. By Theorem 5, k -connected nonhamiltonian graphs must contain an independent set $I = \{x_1, \dots, x_{k+1}\}$ of $k+1$ vertices. The graph is disconnected on removal of the the $n-(k+1)$ vertices of $G-I$, thus we must have $n-(k+1) > k-1$, or $k \leq \frac{n-1}{2}$.

Now consider the independent set I satisfying $\sum_{i=1}^{k+1} d(x_i) = \sigma_{k+1}(G)$. Let the edges in \overline{G} incident on at least one vertex of I be denoted X_I . Then X_I contains $\binom{k+1}{2}$ edges with both endpoints in I and $\sum_{i=1}^{k+1} (n-1-k-d_G(x_i))$ edges with exactly one endpoint in I . Thus, we obtain

$$|E(\overline{G})| \geq |X_I| = \binom{k+1}{2} + (k-1)(n-k-1) - \sigma_{k+1}(G)$$

Using a slight variation of the above result and Lemma 10, [1] also show the following result.

Lemma 13 ([1]). *Suppose $G = cl(G)$ for a nonhamiltonian graph G of order n , and $m \leq \alpha(G)$. Then*

$$|E(\overline{G})| \geq \begin{cases} \frac{m}{2}(n-m) & \text{for } n \text{ odd} \\ \frac{m}{2}(n-m) + \frac{m}{2} - 1 & \text{for } n \text{ even} \end{cases}$$

With the above results, we are ready to proceed to the proof of Theorem 3.

Proof of Theorem 3.

First of all, assume $k \geq 2$ as we already know that when $|E(G)| > \binom{n-1}{2} + 1$, then G is hamiltonian and consequently connected as well. Assume G is non hamiltonian. We may assume $G = cl(G)$, in which case $d(x) + d(y) \leq n-1$ for any two nonadjacent vertices x, y , from Lemma 10. It suffices to prove that $|E(\overline{G})| \geq \binom{n}{2} - ((\binom{n-k}{2} + k^2) = k \cdot n - \frac{3k^2+k}{2}$. Note first that if $\sigma_{k+1}(G) \leq n + k^2 - k - 1$, by Lemma 12

$$|E(\overline{G})| \geq \binom{k+1}{2} + (k+1)(n-k-1) - (n+k^2-k-1) = k \cdot n - \frac{3k^2+k}{2}$$

as desired. We now assume $\sigma_{k+1}(G) \geq n + k^2 - k$ and show that in this case, $|E(\overline{G})|$ is *strictly* greater than $k \cdot n - \frac{3k^2+k}{2}$.

Case 1: Assume $n > \frac{(k^2-1)\alpha+y}{k}$, where $y = \frac{-k^3+4k^2+3k+2}{2}$.

Let $I = \{x_1, x_2, \dots, x_{k+1}\}$ be a set of $k+1$ independent vertices satisfying $\sum_{i=1}^{k+1} d(x_i) = \sigma_{k+1}(G)$, and assume without loss of generality that

$$d(x_1) \geq \frac{\sigma_{k+1}(G)}{k+1} \geq \frac{n+k^2-k}{k+1}$$

Subcase 1a: Suppose $d(x_1) \geq n-2k$. Note that $V(G) - I - N(x_1)$ is non-empty, as otherwise we would have $d(x_1) = n-k-1$, giving $d(x_i) \leq k$ for $2 \leq i \leq k+1$ as $d(x_1) + d(x_i) \leq n-1$ for $2 \leq i \leq k+1$. This contradicts $\sigma_{k+1}(G) \geq n+k^2-k$. Thus, pick some $v \in V(G) - I - N(x_1)$, giving $d_{\overline{G}}(v) = n-1 - d_G(v) \geq d_G(x_1) \geq n-2k$. Therefore, \overline{G} contains at least $n-2k - |I| = n-3k-1$ edges with both endpoints not in I . Using the same bound we got in Lemma 12 but also including the extra edges in \overline{G} incident with v (that have no endpoint in I) and using Theorem 11, we obtain

$$\begin{aligned} |E(\overline{G})| &\geq \binom{k+1}{2} + (k+1)(n-k-1) + (n-3k-1) - \sigma_{k+1}(G) \\ &\geq (k+2) \cdot n - \frac{k^2+9k+4}{2} - (n+(k-1)\alpha-k) \\ &> k \cdot n - \frac{3k^2+k}{2} + \frac{3k^2+k}{2} - \frac{k^2+9k+4}{2} + k + \frac{(k^2-1) \cdot \alpha + y}{k} - (k-1)\alpha \\ &= (k \cdot n - \frac{3k^2+k}{2}) + \frac{(k-1) \cdot \alpha + y + k(k^2-3k-2)}{k} > (k \cdot n - \frac{3k^2+k}{2}) \end{aligned}$$

as desired, where the last inequality follows from $y = \frac{-k^3+4k^2+3k+2}{2}$.

Subcase 1b: Suppose next that $d(x_1) \leq n-2k-1$. Then there exist distinct vertices $v_1, v_2, \dots, v_k \in V(G) - I - N(x_1)$, and \overline{G} contains at least

$$(d_{\overline{G}}(v_1) - k - 1) + (d_{\overline{G}}(v_2) - k - 2) + \dots + (d_{\overline{G}}(v_k) - 2k) = \sum_{i=1}^k d_{\overline{G}}(v_i) - \frac{3k^2+k}{2}$$

edges with neither endpoint in I . Using $d(v_i) + d(x_1) \leq n-1$ as $G = cl(G)$, we get $d_{\overline{G}}(v_i) \geq d_G(x_1) \geq \frac{n+k^2-k}{k+1}$ for all $1 \leq i \leq k$. Consequently, we obtain at least

$$\frac{k(n+k^2-k)}{k+1} - \frac{3k^2+k}{2}$$

edges in \overline{G} with neither endpoint in I . Using Theorem 11 and Lemma 12 again, we get

$$\begin{aligned}
 |E(\overline{G})| &\geq \binom{k+1}{2} + (k+1)(n-k-1) + \frac{k(n+k^2-k)}{k+1} - \frac{3k^2+k}{2} - (n+(k-1)\alpha-k) \\
 &= (kn - \frac{3k^2+k}{2}) + \frac{k}{k+1}(n) - (k-1)\alpha + \binom{k+1}{2} - (k+1)^2 + \frac{k(k^2-k)}{k+1} + k \\
 &> (kn - \frac{3k^2+k}{2}) + \frac{k}{k+1} \frac{(k^2-1)\alpha+y}{k} - (k-1)\alpha + \frac{-k^2-k-2}{2} + \frac{k(k^2-k)}{k+1} \\
 &= (kn - \frac{3k^2+k}{2}) + \frac{1}{k+1} \left(\frac{-k^3+4k^2+3k+2}{2} + \frac{(-k^2-k-2)(k+1)}{2} + k^3 - k^2 \right) \\
 &= kn - \frac{3k^2+k}{2}
 \end{aligned}$$

as desired.

Case 2: Assume $n \leq \frac{(k^2-1)\alpha+y}{k}$.

In this case, $\alpha \geq \frac{nk-y}{k^2-1}$. By Lemma 13, $|E(\overline{G})| \geq \frac{1}{2}\alpha(n-\alpha)$. This is an upward facing parabola for fixed n , so for $\frac{nk-y}{k^2-1} \leq \alpha \leq n - \frac{nk-y}{k^2-1}$, this function is minimized at $\alpha = \frac{nk-y}{k^2-1}$. Therefore, in this range

$$\begin{aligned}
 |E(\overline{G})| &\geq \frac{\alpha}{2}(n-\alpha) \geq \frac{1}{2} \left(\frac{nk-y}{k^2-1} \right) \left(\frac{n(k^2-k-1)+y}{k^2-1} \right) \\
 &= \frac{n^2k(k^2-k-1) + n(2k+1-k^2)y - y^2}{2(k^2-1)^2}
 \end{aligned}$$

If we want the above to be strictly greater than $kn - \frac{3k^2+k}{2}$,

$$\frac{n^2k(k^2-k-1)}{2(k^2-1)^2} \geq kn \iff n \geq \frac{2(k^2-1)^2}{k^2-k-1} = 2(k^2+k + \frac{1-k}{k^2-k-1})$$

suffices. This is because for $k \geq 5$, $y = \frac{-k^3+4k^2+3k+2}{2} < 0$ and $2k+1-k^2 < 0$, giving $(2k+1-k^2)y > 0$. Similarly, $-y^2 = \frac{(-k^3+4k^2+3k+2)^2}{4} > -(3k^2+k)(k^2-1)^2$ for $k \geq 5$, so we only have to check the cases of $k = 2, 3, 4$ manually which is a routine check.

Now, it remains to consider the possibility that $\alpha > n - \frac{nk-y}{k^2-1} = \frac{n(k^2-k-1)+y}{k^2-1}$. In this case however, α is quite large compared to n , so the $\binom{\alpha}{2}$ edges in \overline{G} between the vertices of an independent set of size α is strictly greater than $k \cdot n - \frac{3k^2+k}{2}$ for all n . Indeed, we manually verify for $k \leq 3$, and for $k \geq 4$ simply note that $\frac{nk}{2} + y \geq 0$, and hence when $n \geq 2(k^2+k)$ we have

$$\alpha > \frac{n(k^2 - \frac{3k}{2} - 1)}{k^2-1} \geq \frac{9n}{15}, \quad \binom{9n/15}{2} > \frac{9n}{30} \cdot \frac{8n}{15} > kn$$

We now prove that the extremal nonhamiltonian k -connected graphs are unique for $n \geq 2(k^2+k)$. Recall that we may assume $G = cl(G)$ is a nonhamiltonian, k -connected graph of order $n \geq 2k^2+2k$ with $\sigma_{k+1}(G) = n+k^2-k-1$

as equality only holds if all the inequalities in the above proof are tight.

Thus, all the edges in \overline{G} have at least one endpoint in I . Let $I = \{x_1, x_2, \dots, x_{k+1}\}$ be a set of independent vertices such that $k \leq d(x_1) \leq \dots \leq d(x_{k+1})$. As mentioned in the previous section, we may further assume that all edges in \overline{G} have at least one endpoint in I , that is, if $x, y \in V(G) - I$, then $\{x, y\} \in E(G)$. We will now use the properties of graph closure repeatedly. First, note that we must have a clique on the remaining $n - k - 1$ vertices, each of which has degree at least $n - k - 2$.

- Say $d(x_k) \geq k + 1$ Consider the neighbours of x_k in the clique. These neighbours have degree at least $n - k - 1$, and hence since $G = cl(G)$, must be adjacent to x_{k+1} as well as $d(x_{k+1}) \geq k + 1$. But then, these neighbours have degree at least $n - k$, and hence must be adjacent to all of x_1, \dots, x_{k+1} by the same argument. Thus, I and $N(I)$ together form a complete bipartite graph with $|N(I)| \geq k + 1 = |I|$. If $d(x_{k+1}) > k + 1$, then it is easy to see that the graph is hamiltonian, and otherwise $k + 1 = d(x_i) \forall i \in [k + 1]$, giving

$$\sigma_{k+1} = n + k^2 - k - 1 = (k + 1)^2 \iff n = 3k + 2$$

which is false as we assumed $n \geq 2k^2 + 2k$

- Otherwise $d(x_k) = k$, and hence $d(x_{k+1}) = \sigma_{k+1} - k^2 = n - k - 1$, so we have a clique on the $n - k$ vertices in $G \setminus \{x_1, \dots, x_k\}$. The neighbours of any $x_i, i \in [k]$ must have degree at least $n - k$, and hence are joined to all the x_i . Thus, we obtain the desired extremal graph with exactly $\binom{n-k}{2} + k^2$ many edges, namely a clique on $n - k$ vertices and k other independent vertices forming a complete bipartite graph with some k vertices from the clique.

4 Concluding Remarks

A simpler proof of Theorem 1 with a weaker constant can be obtained using Turán's theorem and a theorem of Erdős and Gallai [18] on the length of the longest cycle in a graph. Consider any k -cyclable graph with $\alpha(G) \geq k$. Then, let S be a set of k independent vertices, and consider the cycle containing it. This gives us a cycle of length at least $2k$, as any two independent vertices are not adjacent to each other. Thus, we must have $\alpha(G) < k$. By a variant of Turán's theorem, we also have $\alpha > \frac{n}{\tilde{d} + 1}$, where \tilde{d} is the average degree. Thus, we obtain

$$\frac{2|E(G)|}{n} + 1 = \tilde{d} + 1 > \frac{n}{\alpha} \geq \frac{n}{k - 1} \implies |E(G)| \geq \frac{1}{2}n \left(\frac{n}{k - 1} - 1 \right)$$

which is larger than $\frac{1}{2}(2k - 1)(n - 1)$ if $n \geq 2k^2$. giving $c(G) \geq 2k$ when $k \leq \sqrt{n/2}$

It is also interesting to understand what happens to the circumference of k -cyclable graphs for large values of k . As mentioned earlier in the introduction,

it is not necessarily the case that $c(G) = n$ when $k = n - 1$ due to the existence of hypohamiltonian graphs. Thus, we have the following extremal problem.

Conjecture 1. For a given n , let $f(n)$ be the largest value of k such that any k -cyclable graph satisfies $c(G) > k$. From the above, we have $f(n) < n - 1$ and from Theorem 1, we know $f(n) = \Omega(n)$. Is it the case that $f(n) = n - 2$?

We can also ask for what regime of k as a function of n do results of the type in Theorem 1 hold.

Conjecture 2. For a given n , let $g(n)$ be the largest value of k such that any k -cyclable graph satisfies $c(G) \geq 2k$. From Theorem 1 we know $g(n) = \Omega(\sqrt{n})$. Is it the case that $g(n) = O(\sqrt{n})$?

Moreover, our results only give an improvement of the form $c(G) \geq (1 + \gamma)k$, $0 < \gamma < 1$, for k up to around $2\sqrt{n}$, and it is natural to ask if such a linear bound on the circumference can be obtained for much larger regimes of k . Finally, note that the results of Theorem 3 only hold for $n \geq 2(k^2 + k)$. For fixed values of $k \leq 3$, [1] give a tight bound for the minimum value of n for this to hold. They also note that this bound cannot hold for $k = \Omega(n)$, in particular the graph obtained by joining $n - k$ independent vertices to each vertices of K_k is k -connected and non hamiltonian, with total number of edges more than $\binom{n-k}{2} + k^2$ when $\frac{n+1}{6} < k < \lfloor \frac{n-1}{2} \rfloor$. This still leaves a significant gap in the possible range of k for which k -connectivity and $|E(G)| > \binom{n-k}{2} + k^2$ implies hamiltonicity, as our result only applies for $k = O(\sqrt{n})$.

References

1. Owen D. Byer, Deirdre L. Smeltze, Edge bounds in nonhamiltonian k -connected graphs, *Discrete Mathematics* 307 (2007) 1572–1579
2. G.A. Dirac, In abstrakten graphen vorhandene vollständige 4-graphen und ihre unterteilungen, *Math. Nachr.* 22 (1960) 61–85,
3. R.J. Gould, A look at cycles containing specified elements of a graph, *Discrete Math.* 309 (21) (2009) 6299–6311,
4. G. Dirac, Some theorems on abstract graphs, *Proceedings of the London Mathematical Society* 2 (1952) 69–81.
5. V. Chvátal, P. Erdős, A note on hamiltonian circuits, *Discrete Math.* 2 (1972) 111–113.
6. Bauer, L. McGuire, H. Trommel, H.J. Veldman, Long cycles in 3-cyclable graphs, *Discrete Mathematics* 218 (2000) 1–8
7. A. Björklund, T. Husfeldt, N. Taslamán, Shortest cycle through specified elements, in: *Proceedings of the Twenty-Third Annual ACM-SIAM Symposium on Discrete Algorithms, SODA 2012, Kyoto, Japan, January (2012)* 17-19, SIAM, 2012, pp. 1747–1753.
8. Christophe Crespelle, Petr A. Golovach, Cyclability in graph classes, *Discrete Applied Mathematics* 313 (2022) 147–178

9. K. Kawarabayashi, An improved algorithm for finding cycles through elements, *Integer Programming and Combinatorial Optimization*, 13th International Conference, IPCO 2008 (Bertinoro, Italy, May 26–28, 2008), Proceedings. Lecture Notes in Computer Science 5035, Springer 2008, pp. 374–384.
10. J. Doyen, V.V. Diest, (1975), New families of hypohamiltonian graphs, *Discrete Mathematics*, 13 (3): 225–236
11. L. Ng, M. Schultz, k -Ordered Hamiltonian graphs, *J. Graph Theory* 1 (1997) 45–57.
12. Ralph J. Faudree, Survey of results on k -ordered graphs, *Discrete Mathematics* 229 (2001) 73–87
13. J.L. Fouquet, J.L. Jolivet, Problème 438, in: *Problèmes combinatoires et théorie des graphes*, Univ. Orsay, Orsay, 1976.
14. Suil O, Douglas B. West, Hehui Wu, Longest cycles in k -connected graphs with given independence number, *Journal of Combinatorial Theory, Series B*, Volume 101, Issue 6, November 2011, Pages 480–485
15. J.A. Bondy, V. Chvátal, A method in graph theory, *Discrete Math.* 15 (1976) 111–135.
16. O. Ore, A note on hamiltonian circuits, *Amer. Math. Monthly* 67 (1960) 55
17. Hao Li, Generalizations of Dirac’s theorem in Hamiltonian graph theory—A survey, *Discrete Mathematics* 313 (2013) 2034–2053
18. P. Erdős and T. Gallai, On maximal paths and circuits of graphs, *Acta Math. Hungar.*, 10 (1959), 337–356.